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A SEQUENTIAL k-GROUP RANDOM ALLOCATION METHOD WITH APPLICATIONS TO SIMULATION

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A Sequential k-Group Random Allocation Method with Applications to Simulation

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A sequential method of random allocation is given and it is shown how it can be used to estimate the observed significance levels of k-sample nonparametric tests. The sequential technique is compared to the standard random allocation technique and shown to be more efficient. An application is made to the Dunn-Bonferroni method of multiple comparisons.

AMS (MOS) Subject Classification: 621.99, 65C05

Key Words: Dunn-Bonferroni method; Nonparametric tests; Observed significance levels; Simulation

Work Unit No. 4 - Statistics and Probability

1. The Sequential Allocation Method

Bebbington (1975) showed that if there were N objects (such as file cards) from which it was desired to select (without replacement here and throughout) a random sample of size k without numbering the N objects, then one could proceed sequentially by selecting the first object with probability k/N and if at the T^{th} stage s have been selected, then the $T^{+|\frac{nt}{n}|}$ object is selected with probability (k-s)/(N-T), $T = 1, 2, \ldots, N-1$.

We now state and prove the extension to an arbitrary number of groups. Suppose there are N objects and it is desired to sequentially divide them randomly into r groups of size k_1,k_2,\ldots,k_r , $\sum_{i=1}^{L}k_i=N,\ i.e.,\ \text{each allocation has probability}$ of $1/\left\{k_1,\ldots,k_r\right\}.$ Let s_1r,\ldots,s_{rT} be the number of objects selected for groups $1,2,\ldots,r$ at the $T^{\underline{th}}$ stage and let P_i,r_i denote the selection probability for group i at the T^{+1} stage. Then if

$$P_{1,T+1} = (k_1 - s_{1T})/(N-T), T = 0,1,...,N-1, (1.1)$$

the selection is random. Note that $P_{1,1}=k_1/N$ and $\sum\limits_{i=1}^{E}P_{1,T+1}=1$. The randomness follows immediately by noting that the probability of a particular assignment is

$$\binom{E}{1-1}k_1$$
: $\binom{E}{1-1}(k_1, \dots, k_E)$.

Bebbington's (1975) result is a special case of the above when $r=2\,.$

As an example, suppose r=3, $k_1=2$, $k_2=2$, $k_3=3$ and N=7. In order to make the sequential allocation given by (1.1) we take

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? independent random numbers U_1 , i=1,2,...,7. Let

$$Q_{0,T} = 0$$
 and $Q_{1,T} = \sum_{j=1}^{1} P_{j,T}$, i = 1,2,...,r, T = 1,2,...,N.

Then the m^{th} object, $m=1,2,\ldots,N$, is assigned to group n, where n is the unique integer such that

Suppose the 7 random numbers are .79039, .01850, .99744, .81812, .93169, .22705, and .97709. The selection process is summarized in Table 1.

Insert Table 1 here.

Mote that if all the $\mathbf{k}_{\underline{1}}$'s are one, a random permutation is produced if we think of the group as denoting position.

2. Applications to Simulation

In k-sample nonparametric tests the observed significance level of the test is obtained by considering all possible partitions M of the (possibly tied) observed values or (possibly average) ranks into r groups, computing the value of the test statistic, and counting the number of times m it is equal to or greater than the observed value. The observed significance level â is then m/M. When the number of partitions is large this is prohibitive and â is estimated either by simulation (taking a large random sample of the allocations) or by asymptotics. The advantage of simulation is that one can control the accuracy of the estimate (by taking a large or small random sample) depending on the importance of the situation, unlike asymptotics which each time it is used forces one into the straight-jacket of committing a usually unknown error. Since it is (perhaps regretably) a well known fact that different actions will be taken for close values

of \hat{a} , one above and the other below some fixed level (e.g., .01, .05, or .1) of the decision-maker, the use of simulation at least prevents approximating error in \hat{a} to be the determining factor.

if it is decided to use simulation, then a possible procedure is to make the random assignment as described in Section 1 many times by using a computer. The commonly used method is to produce a random permutation by ordering a random sample of uniform numbers and choosing the first k_1 indexes for group 1, the next k_2 for group 2, and so on. If all the $k_1^{'}$'s are one, then this is more efficient than Section 1. However, as soon as the $k_1^{'}$'s depart even moderately from 1, the method of Section 1 becomes much more efficient. As an example, if $k_1=k_2=k_3=k_4=10$ and it is desired to make 2000 random assignments using a UNIVAC 1110 computer, a FORTRAN program using the methods of Section 1 uses 4.71 seconds of takes 9.17 seconds.

The Appendix contains a listing of the FORTRAN subroutine RANDM that uses the theory of Section 1 to make random assignments. This may be tied in with any specific simulation problem, e.g., the case treated in Section 3.

Applications to the Dunn-Bonferroni Method of Multiple Comparisons

The D-B (Dunn-Bonferroni) method is described in Dunn (1964). Briefly, let Y_{ij} , $i=1,2,\ldots,r$, $j=1,2,\ldots,n_1$, be continuous (this assumption is not important and is removed later) random variables with distribution function F_i , H_0 : $F_1 = F_2 = \ldots = F_r$, H_a : for at least one pair $\{i,j\}$, $F_i \neq F_j$ in the sense of producing larger or smaller values), and the test must identify which, if any, pairs

$$\mathbf{z}_{1,j} = |\vec{\mathbf{z}}_1 - \vec{\bar{\mathbf{z}}}_3| \sqrt{\left[\frac{(n)(n+1)}{12} \left(\frac{1}{n_1} + \frac{1}{n_j}\right)\right]^{1/2}} \ge \mathbf{z}_0/(k(k-1)) \ , \ (3.1)$$

where \overline{k}_1 denotes the average of the ranks of the $i\frac{th}{h}$ group in the joint ranking. The nominal significance level of this procedure is a. The actual significance level α_k is

$$\alpha_{A} = P_{0} \left[\max_{1 \le j} z_{1} / (k(k-1)) \right],$$
 (3.2)

and may be obtained by simulation based on Section 1. Table 2 gives some comparisons of nominal with actual, using Section 1 and 10,000 simulations.

Insert Table 2 here.

The Appendix contains a listing of the program used for Table 2. It thus appears that D-B is conservative and we can remove the conservatism by substituting for $\mathbf{x}_{\alpha}/(k(k-1))$ d_{1}, where d_{{1}}, is 1-1,2,...,r(r-1)/2, is the ith largest observed values of \mathbf{z}_{ij} , if \mathbf{z}_{ij} , to obtain by simulation the r(r-1)/2 possible observed significance levels.

The K-S (Kruskal-Scheffé) method is also sometimes used in this situation (see, e.g., Miller, 1966, p. 166) and consists of replacing $z_{\alpha}/(r(r-1))$ in (3.1) with $h_{\alpha}^{1/2}=(\chi_{\alpha,r-1}^2)^{1/2}$, where $\chi_{\alpha,r-1}^2$ is the upper $\alpha^{\pm h}_{1}$ point of χ^2 with r-1 degrees of freedom. The comparison of the critical constants in Table 3 shows that this is even more conservative than D-B.

Insert Table 3 here.

If the data is discrete, the D-B method can be modified as in Dunn (1964) and the random assignment done on average ranks. Thus ties present no problems in this approach.

The third method discussed in Miller (1964), the Steel manyone rank statistics, is too time-consuming for the simulation approach. For all practical purposes the exact D-B (use of the $d_{(1)}$ and simulation) seems the best method to use.

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New York: McGraw-Hill Book Company.

1. Selection Process

Group		~		n	~		~
P 3T	3,1	3/6	2/2	*	•	•	•
P. T.	1/2	9/2	2/5	2/4	2/3	1/2	-
P 1.T	1/2	3/6	1/5	*	1/3	1/2	•
Random	. 79039	.01850	.99744	. 81812	.93169	. 22705	.97709
Stage	1	~	m	•	v	•	•

Comparison of Actual to Nominal a

	2. Comparison o	2. Comparison of Actual to Nominal a	8		3. Compariso	3. Comparison of D-B and K-S Critical Constants	-S Critical Cor	stants
	Comon	Noninal a	Actual G	14	z.05/(r(r-1))	$z_{.05}/(r(r-1)) \frac{(\chi_{.05;r-1}^2)^{1/2}}{(\chi_{.05;r-1}^2)^{1/2}} \frac{z_{.01}/(r(r-1))}{z_{.01}/(r(r-1))} \frac{(\chi_{.01;r-1}^2)^{1/2}}{(\chi_{.01;r-1}^2)^{1/2}}$	z.01/(r(r-1))	(x,01,r-1)1/2
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,	, 5	50.	.040	→	2.50	3.08	3.02	3.65
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